

Problem 1: Translation. Let \mathcal{L} be a formal language containing the individual constants $a, c_1, c_2, d, e, f, m,$ and $r,$ and the predicates $H, J,$ and $L,$ with the following key for the translation between \mathcal{L} and English:

a : Ashleigh	e : Evan	Px : x was (ever) president of a house
c_1 : Cappie	f : Frannie	Zx : x was (ever) a member of ZBZ
c_2 : Casey	m : Max	Lxy : x fell in love with y
d : Dale	r : Rebecca	

Translate the following English sentences to WFFs of \mathcal{L} as directly and plainly as possible. Remember that not all subtleties of meaning present in an English sentence can be represented yet in our formal logical system.

- Frannie was never a member of ZBZ, and neither was Dale.
- Ashleigh and Evan were each president of a house if Cappie ever was.
- Max was a member of ZBZ if and only if Casey was, too.
- Neither Max nor Rebecca were ever president of a house, although Frannie and Casey both were.
- Dale fell in love with exactly one of Ashleigh and Casey.
- Frannie and Evan fell in love with each other, as did Rebecca and Cappie.

Problem 2: Valuation in a Model. Assume \mathcal{L} as in Problem 1. Further assume a model \mathbb{M} of \mathcal{L} consisting of a domain of entities \mathcal{D} (defined below) and an interpretation function I (partly defined below).

$$\mathcal{D} = \left\{ \begin{array}{l} \text{ASHLEIGH, CAPPIE, CASEY,} \\ \text{DALE, EVAN, FRANNIE,} \\ \text{MAX, REBECCA} \end{array} \right\} \quad \begin{array}{ll} I(a) = \text{ASHLEIGH} & I(e) = \text{EVAN} \\ I(c_1) = \text{CAPPIE} & I(f) = \text{FRANNIE} \\ I(c_2) = \text{CASEY} & I(m) = \text{MAX} \\ I(d) = \text{DALE} & I(r) = \text{REBECCA} \end{array}$$

- Finish defining \mathbb{M} by providing definitions of $I(P), I(Z),$ and $I(L)$ so that sentences 1b, d, and e are all true in \mathbb{M} , and 1a, c, and f are all false in \mathbb{M} . There is more than one possible solution; you can just use the simplest one, if you prefer.
- Show explicitly, step-by-step, how to calculate $V_{\mathbb{M}}(1b) = 1$.
- Show explicitly, step-by-step, how to calculate $V_{\mathbb{M}}(1f) = 0$.

Problem 3: Proof. Symmetric predicates are two-place predicates like *next to*, *married to*, and *touched*, in which the two participants in the predicate can be swapped without changing the truth value. For example, if *Alice is next to Bob* is true, then we expect *Bob is next to Alice* to also be true. It is simply impossible given our usual interpretation of *next to* for two people to be arranged in such a way that one is next to the other, but not vice versa. We will formally define a symmetric predicate as follows:

Definition: Let P be a two-place predicate in a formal language \mathcal{L} , and let \mathbb{M} be a model with a domain of entities \mathcal{D} and an interpretation function I . Then P is symmetric in \mathbb{M} iff for any arbitrary entities $\varepsilon_1, \varepsilon_2 \in \mathcal{D}$, if $\langle \varepsilon_1, \varepsilon_2 \rangle \in I(P)$, then $\langle \varepsilon_2, \varepsilon_1 \rangle \in I(P)$.

Let Pab and Pba be WFFs of a formal language \mathcal{L} , and let \mathbb{M} be a model of \mathcal{L} with a domain of entities $\mathcal{D} = \{AA, BB\}$ and an interpretation function I such that $I(a) = AA$ and $I(b) = BB$. Prove that if P is symmetric in \mathbb{M} , then $V_{\mathbb{M}}(Pab \leftrightarrow Pba) = 1$.